# CS 696 Intro to Big Data: Tools and Methods Fall Semester, 2016 Doc 9 Parallel Computing Sep 22, 2016 

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## Parallel Computing

Concurrent computing

Running multiple processes or threads on same processor

Processes or threads are time-sliced

Parallel computing

Running multiple processes on different processors

Processes in same program run at the same time on different processors

Scaling up

Adding more resources to a machine to allow it to handle larger tasks Memory
Disk space
Faster processor

Scaling out

Adding more machines/processors to handle larger tasks

Requires parallel programming

## Julia Parallel Processing

Low level constructs

High level constructs

Runs on
Multicore processors
Clusters

Cluster management

Experimental Julia to C/C++ compilers from Intel Labs
Run Julia code 20 to 100 times faster than Spark
Spark claims to be 10 to 100 times faster than Hadoop

## Julia Parallel Processing - Low Level

@spawn, @spawnat
Run code on separate/remote processor
@everywhere
Run a command on all processors
fetch
Obtain results from separate processor
put!
store a value on a separate processor


## Adding 1 elementwise In parallel

addprocs(2)
workers()
procs()
remote $=$ @spawn rand(2,2)
fetch(remote)
result = @spawn 1 .+ fetch(remote)
fetch(result)
\# only have two cores
\# [ 2, 3]
\# [ 1, 2, 3]
\# RemoteRef\{Channel\{Any\}\}(2,1,3)
\#= [0.477549 0.193374;
$0.2507990 .0512077]$
\#[1.47755 1.19337;
$1.25081 .05121]$
$\mathrm{A}=\mathrm{rand}(\mathrm{n}, \mathrm{n})$
Aref = @spawn sum( $\mathrm{A}^{\wedge} 2$ ) fetch(Aref)

Bref = @spawn sum(rand(n,n)^2) fetch(Bref)


You need to be aware what you are doing on each processor

Do you need to send A from P1 to P2?

```
function count_heads(n)
    c: : Int = 0
    for \(\mathrm{i}=1\) : n
        c += rand(Bool)
    end
    C
end
a = @spawn count_heads(100000000) On worker 5:
b = @spawn count_heads(100000000) function count_heads not defined on process 3
fetch(a)+fetch(b)
```

count_heads.jl
function count_heads(n)
c::Int $=0$
for $i=1: n$
$c+=$ rand(Bool)
end
c
end

Put count_heads.jl in Julia path

In Julia

```
require("count_heads")
a = @spawn count_heads(100000000)
b = @spawn count_heads(100000000)
fetch(a)+fetch(b)
```

High-level Parallel/Performance Constructs
@parallel

$$
\begin{aligned}
& \text { @parallel reducer for var = range } \\
& \text { body } \\
& \text { end }
\end{aligned}
$$

Divide the loop among worker processes

Each process accumulates results and used reducer to combine the results

Result is send back to master and reduce is used combine all results

```
addprocs(10)
@parallel (+) for k = 1:100_000
    rand(1)
end
```

Each worker will sum 10_000 random numbers

Master will sum up the 10 results

Assuming you have 10 processors

## Computing Pi




Area $=\mathrm{pi} / 4$


Area of Square $=1$

Select random point in unit square

Probability that point is in the quarter Circle is (pi/4)/1 = pi/4

Select N random points in unit square
Let $\mathrm{K}=$ number of points in quarter circle
K should be about $\mathrm{N}^{*} \mathrm{pi} / 4$
$4 \mathrm{~K} / \mathrm{N}$ should be about pi

## Computing Pi

| rand(2) |  |  |
| :---: | :---: | :---: |
| function findpi(n) inside $=0$ | returns two random numbers between 0 and 1 |  |
| for $\mathrm{i}=1: n$ |  |  |
| $x, y=\operatorname{rand}(2)$ |  |  |
| $\begin{aligned} & \text { if }\left(x^{\wedge} 2+y^{\wedge} 2<=1\right) \\ & \text { inside }+=1 \end{aligned}$ |  |  |
| end | N | findpi(N) |
| end | I_000 | 3.148 |
| 4 *inside / n | 100_000 | 3.15028 |
|  | 100_000_000 | 3.14169832 |
|  | I_000_000_000 | 3.141595912 |

$$
\pi=3.1415926535897 \ldots
$$

## Parallel Version

```
function parallel_findpi(n)
    inside = @parallel (+) for \(\mathrm{i}=1: \mathrm{n}\)
        \(x, y=r a n d(2)\)
        \(x^{\wedge} 2+y^{\wedge} 2<=1 ? 1: 0\)
    end
    4 * inside / n
end
```


## findpi verses parallel_findpi

On Edora - Has 4 Intel(R) Xeon(R) CPU E5-2680 v3 @ 2.50GHz processors addprocs(4)

Time to Run in Seconds

| N | findpi(N) | parallel_findpi(N) | Speedup |
| ---: | :---: | :---: | :---: |
| I_000 | 0.000079 | 0.003326 | 0.02 |
| $100 \_000$ | 0.012105 | 0.06380 | 0.19 |
| 100_000_000 | 7.234 | 2.19617 | 3.29 |
| I_000_000_000 | 72.36 | 18.660 | 3.88 |

## Speedup

$T(1)$ = time for sequnetial program to run
$T(N)=$ time for parallel program to run on $N$ processors
$\mathrm{S}(\mathrm{N})=$ speedup using N processors

$$
S(N)=T(1) / T(N)
$$

## Timings on JuliaBox - 16 CPU

Time to Run in Seconds

| N | findpi(N) | parallel_findpi(N) <br> I6 processor | parallel_findpi(N) <br> 8 processor |
| ---: | ---: | ---: | ---: |
| $100 \_000$ | 0.009 | 0.009 | 0.010 |
| I_000_000 | 0.083 | 0.030 | 0.069 |
| $10 \_000 \_000$ | 0.813 | 0.234 | 0.216 |
| I00_000_000 | 8.143 | I .656 | 2.133 |
| I_000_000_000 | 82.219 | 14.663 | 20.762 |
| $10 \_000 \_000 \_000$ |  | 125.585 | 246.524 |

Speedup

| $100 \_000$ | 0.009 | 1 | 0.9 |
| ---: | ---: | ---: | ---: |
| I_000_000 | 0.083 | 2.8 | 1.2 |
| $10 \_000 \_000$ | 0.813 | 3.5 | 3.8 |
| $100 \_000 \_000$ | 8.143 | 4.9 | 3.8 |
| I_000_000_000 | 82.219 | 5.6 | 4.0 |
| $10 \_000 \_000 \_000$ |  | 6.5 | 3.3 |

## Amdahl's Law

Ts = time of task that is inherently sequential

Tp = time of task that can be parallelized
$T(1)$ = time for sequnetial program to run
$T(N)=$ time for parallel program to run on $N$ processors
$S(N)=$ speedup using $N$ processors

$$
T(1)=T s+T p
$$

$T(N)=T s+T p / N \quad$ Assuming we can parallelize perfectly

$$
\begin{aligned}
S(N) & =T(1) / T(N) \\
& =(T s+T p) /(T s+T p / N)
\end{aligned}
$$

## Amdahl's Law

$S(N)=(T s+T p) /(T s+T p / N)$
if Ts $=0$ and we can perfectly parallelize the task we get
$\mathrm{S}(\mathrm{N})=\mathrm{Tp} /(\mathrm{Tp} / \mathrm{N})=\mathrm{N}$

Ts is never zero
Perfect parallization is not possible

So

$$
S(N)<N
$$

## Amdahl's Law

Theory

$$
\begin{aligned}
& S(N)=(T s+T p) /(T s+T p / N) \\
& S(N)<N
\end{aligned}
$$

## Practice

It is possible for $S(N)>N$

How
Single processor may not be able to fit data in physical memory Paging will significally slow sequenial program down N processors can have more total memory that single processor So parallel version may not have paging issues

## More Realistic Amdahl's Law

Ts = time of task that is inherently sequential

Tp = time of task that can be parallelized

Tis = Average additional serial time doing interprocessor communication
Assume each processor takes same amount of time
Total time is $\mathrm{N}^{*}$ Tis

Tip - Average additional time by each processor doing set up, idle time, etc.
$S(N)=(T s+T p) /\left(T s+N^{*} T i s+T i p+T p / N\right)$
$S(N)=(T s+T p) /\left(T s+N^{*} T i s+T i p+T p / N\right)$

Ts = Tip $=0$
Tp = 10_000
Tis $=20=0.2 \%$ * $T p$


## Monte Carlo Method

Uses repeated random sampling to obtain numerical results

Used mainly in:

Optimization<br>Numerical Integration<br>Generating draws from probability distribution

## Embarrassingly (Pleasingly) Parallel

Little or no effort needed to separate problem into parallel tasks

Little or no communication needed between parallel tasks

Searching a web page that contains key words

These are the types of problems that can be

Solved using Hadoop \& Spark

Compilers can detect some forms and parallelize for you

## Distributed Arrays - DistributedArrays.jI

Distributes Arrays among processors

Can distribute arrays from master to slaves

Can create arrays on slaves

Master can work arrays on slaves

Distributing Data


## Creating a Distributed Array

```
dzeros(100,100,10)
dones(100,100,10)
drand(100,100,10)
drandn(100,100,10)
dfill(x,100,100,10)
```


## Using DistributedArrays.jl

onmaster $=$ rand $(100,100)$
distributed $=$ distribute(onmaster)
sum(distributed)
heads $=\operatorname{map}(x->x>0.5$, distributed $)$
\#distribute onmaster to the workers
\# compute sum locally on workers \# combine the result on master
\# apply map on workers
\$ return result on master

## SharedArrays

Each worker has access to the array

```
addprocs(3)
    3-element Array{Int64,1}:
    2
    3
    4
S = SharedArray(Int, (3,4), init = S -> S[Base.localindexes(S)] = myid())
```

    \(3 \times 4\) SharedArray\{Int64,2\}:
    2234
    2334
    2344
    
## ClusterManagers

Launches worker processes in a cluster environment
Managing events during the lifetime of each worker
Providing data transport

Julia Cluster

The initial Julia process, also called the master, is special and has an id of 1
Only the master process can add or remove worker processes
All processes can directly communicate with each other

## Types of Cluster Managers

LocalManager,
used when addprocs() or addprocs(np::Integer) are called

SSHManager
used when addprocs(hostnames::Array) is called with a list of hostnames

Remote hosts need passwordless login enabled

## ArrayFire.jl

GPU computing
using ArrayFire
$a=\operatorname{rand}(10,10)$
on_gpu = AFArray(a)
result_on_gpu $=($ on_gpu +1$) / 5$
result_on_cpu = Array(result_on_cpu)

## HPAT.jl, ParalleIAccelerator.jl

Intel Labs projects to provide high level efficient \& fast parallel code

ParallelAccelerator.jl
Converts Julia code to C/C++ Imports C/C++ code into Julia

Supports subset of Julia

Uses implicit parallelism in map, reduce, comprehension
.+, .- , .* , ./ converted into data-parallel map operations

HPAT.jI
Using ParallelAccelerator converts Julia code to
C/C++ \& MPI calls for distributed computing

## Sample Using ParallelAccelertor


function calc_pi_normal(n)
$x=\operatorname{rand}(n)$.* 2.0 - 1.0
$\mathrm{y}=\operatorname{rand}(\mathrm{n})$.* 2.0 . 1.0
return $4.0^{*}$ sum(x.^2 .+ y. ${ }^{\wedge} 2 .<$
end
@time calc_pi(10_000_000)
0.284697 seconds
(28 allocations: 1.641 KB )
@time calc_pi_normal(10_000_000)
1.167740 seconds
(7.49 k allocations: 688.223 MB, $52.57 \%$ gc time)

Using for loop rather than .* etc
1.105030 seconds
(10.00 M allocations: 915.528 MB, $18.67 \%$ gc time)

## Sample Using HPAT

using HPAT

@acc hpat function calc_pi(n)
$x=\operatorname{rand}(n)$.* 2.0 .- 1.0
$y=\operatorname{rand}(n)$.* 2.0 .- 1.0
return 4.0 * $\operatorname{sum}\left(x .{ }^{\wedge} 2 .+y . \wedge 2 .<1.0\right) / n$
end

Now can be run on machines supporting mpi

## HPAT vs. Spark

- Spark ■HPAT


Cori at NERSC/LBL 64 nodes (2048 cores)

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